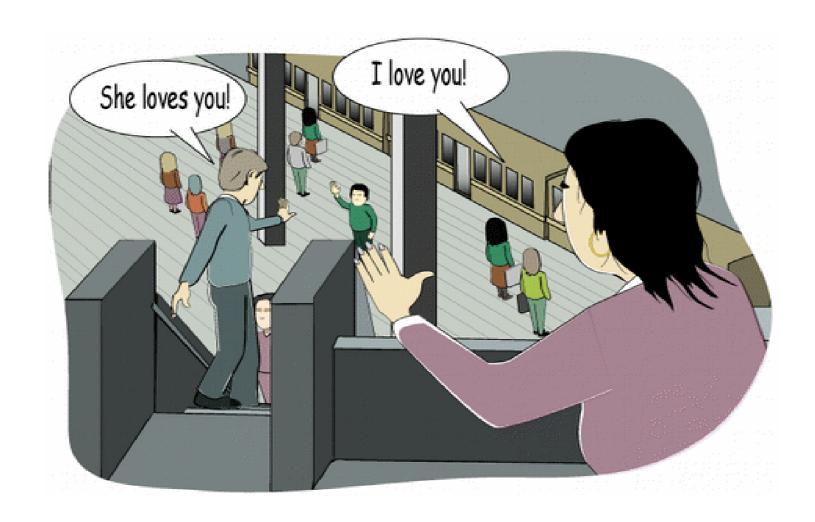
Cybersecurity TLS

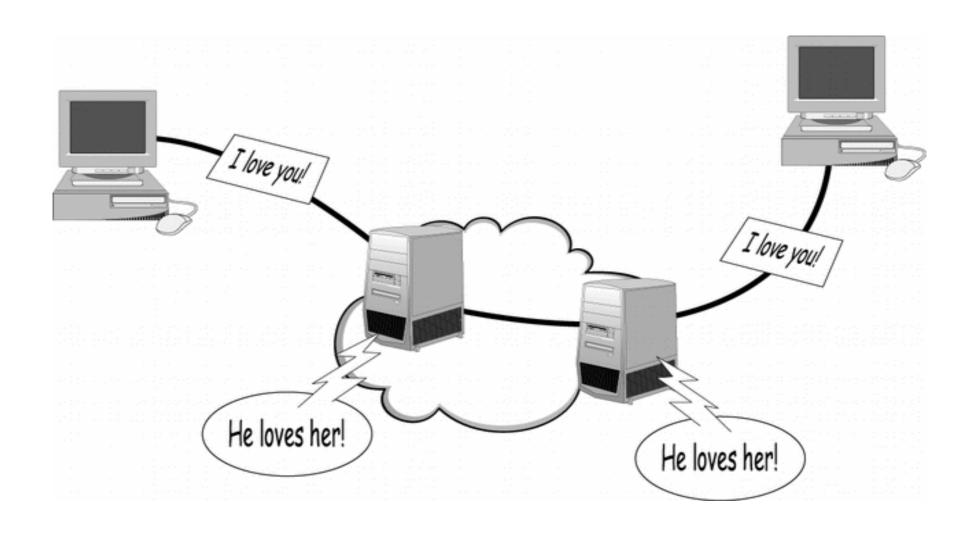
Kasun De Zoysa

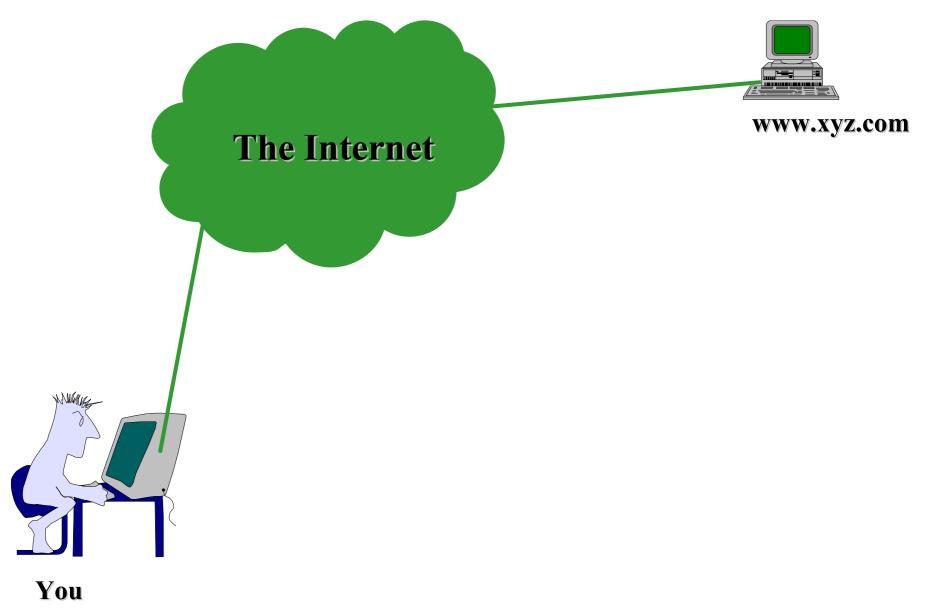
Department of Communication and Media Technologies
University of Colombo School of Computing
University of Colombo
Sri Lanka

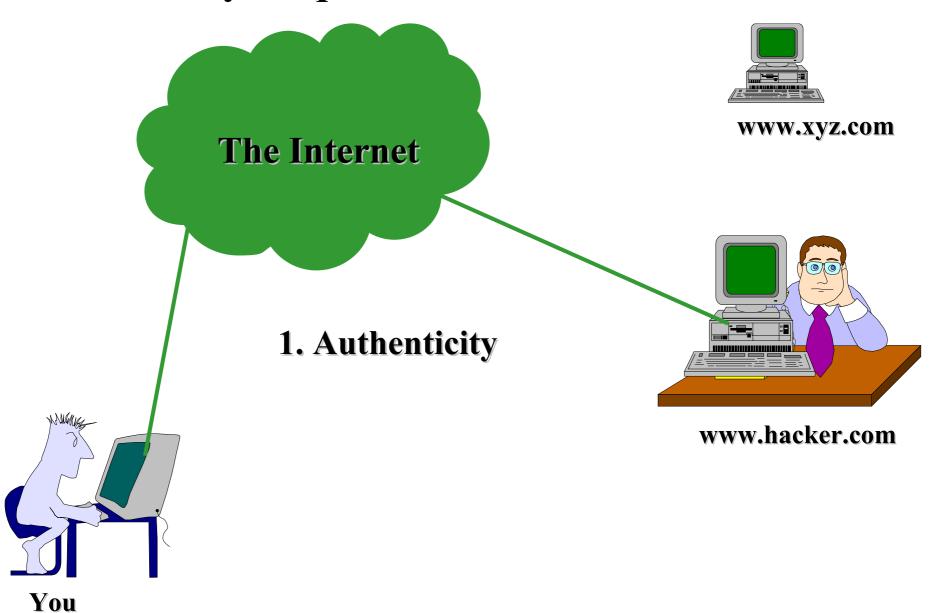
How the Internet Works -1

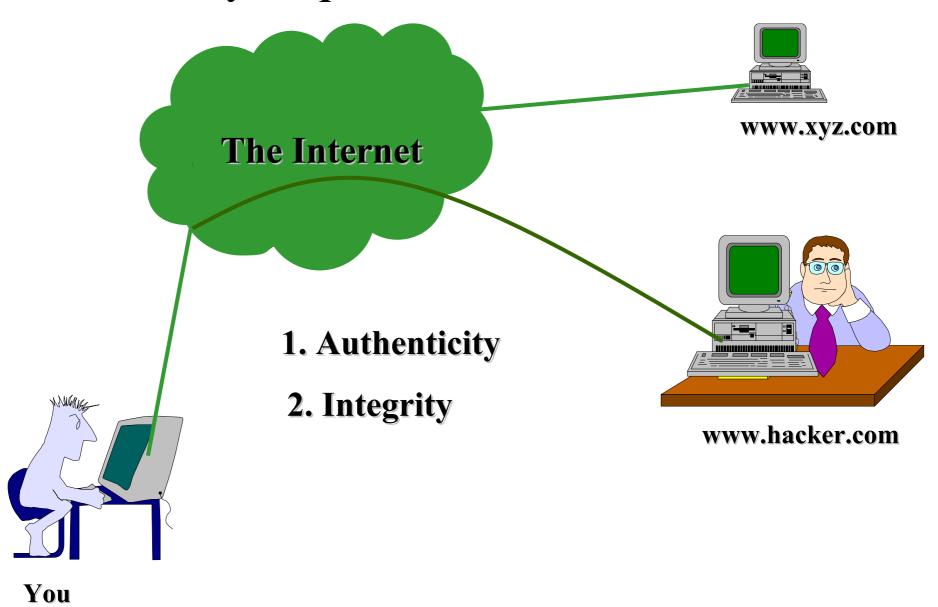


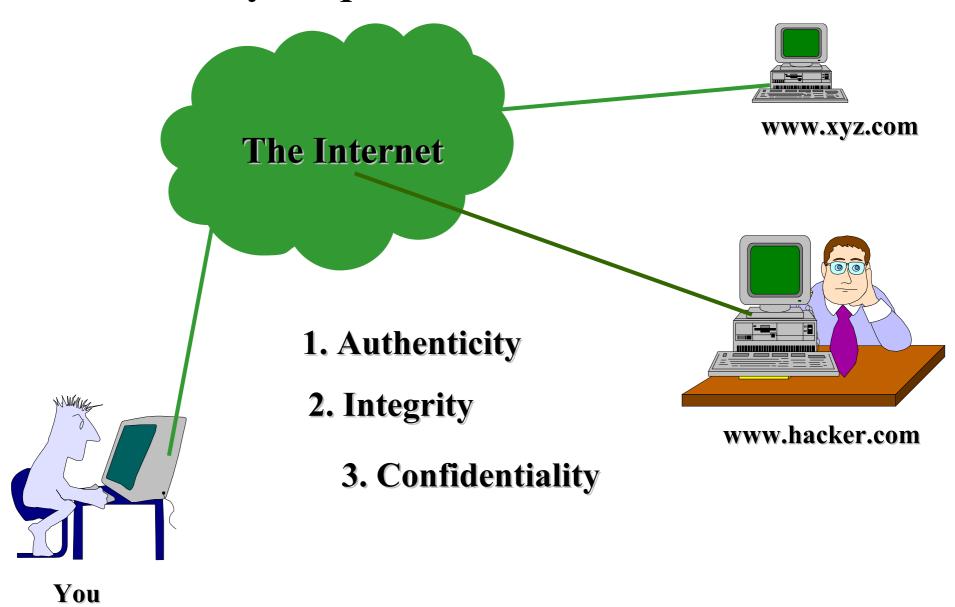
How the Internet Works -2

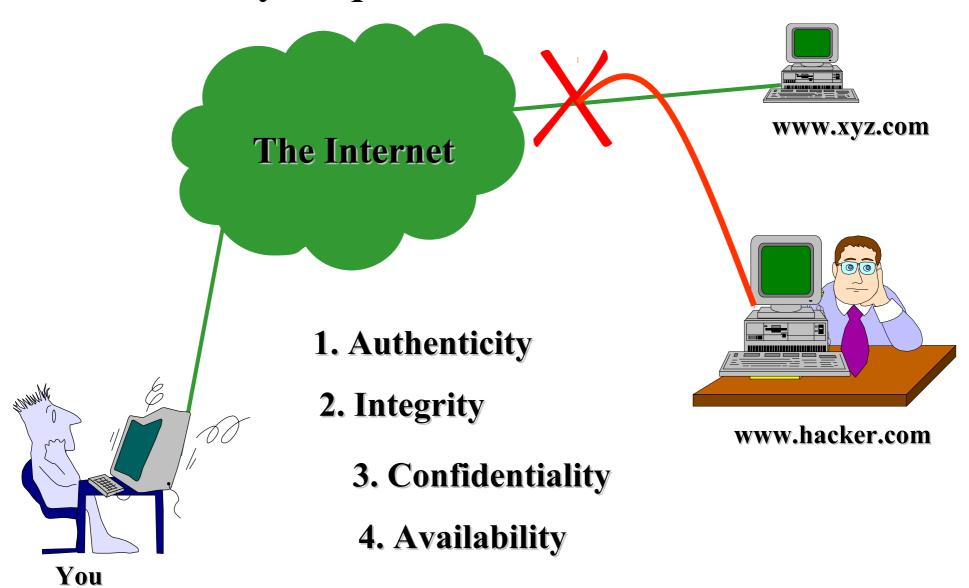


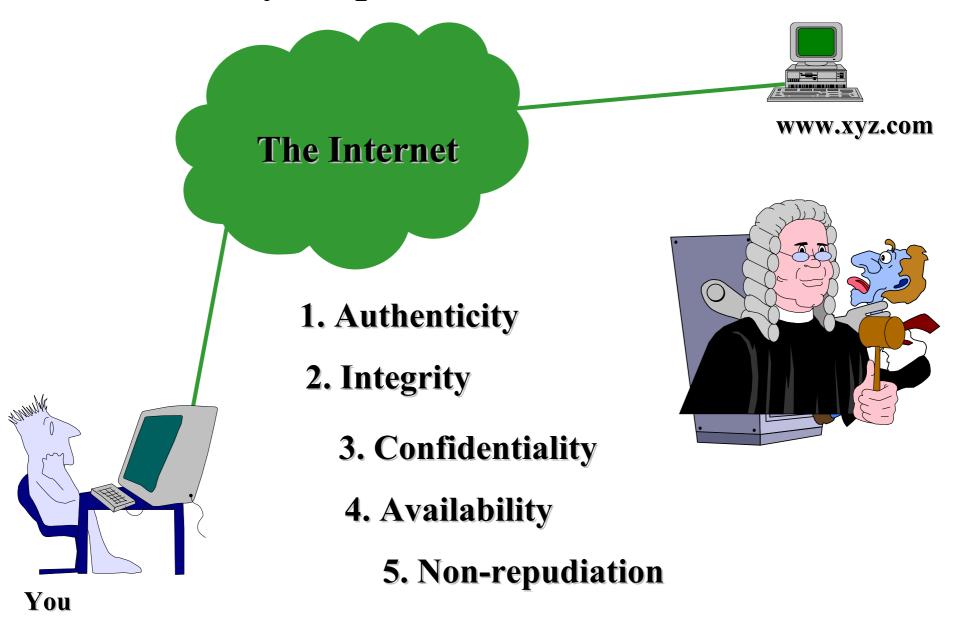












Solutions

Protection at Two Levels:

- 1. Lower Level (Channel protection) (Communication security services)
- 2. Application/User Level (Application level security services)

Secure Socket Layer History

- SSL 1.0 Netscape 1994
- S-HTTP (web only)
- SSL 2.0 Netscape (buggy)
- PCT Microsoft (loser) 1996
- SSL 3.0 Netscape
- TLS 1.0 IETF 1999
- TLS 1.2 now dominant

TLS: Transport Layer Security

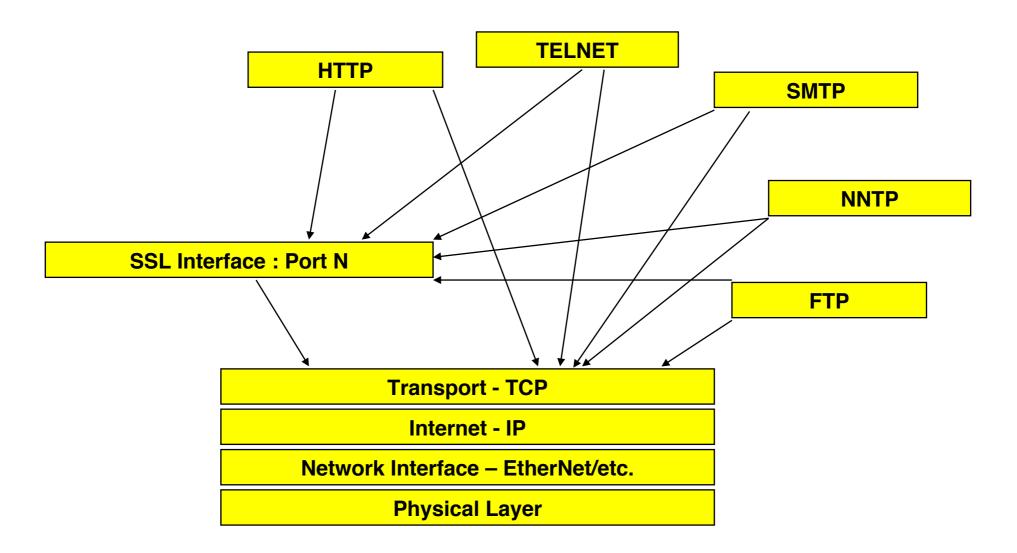
- formerly known as
 - SSL: Secure Sockets Layer
- Addresses issues of privacy, integrity and authentication
 - What is it?
 - How does it address the issues?
 - How is it used

What is TLS?

- Protocol layer
- Requires reliable transport layer (e.g. TCP)
- Supports any application protocols

HTTP	Telnet	FTP	LDAP	
TLS				
TCP				
IP				

Protocol Stack

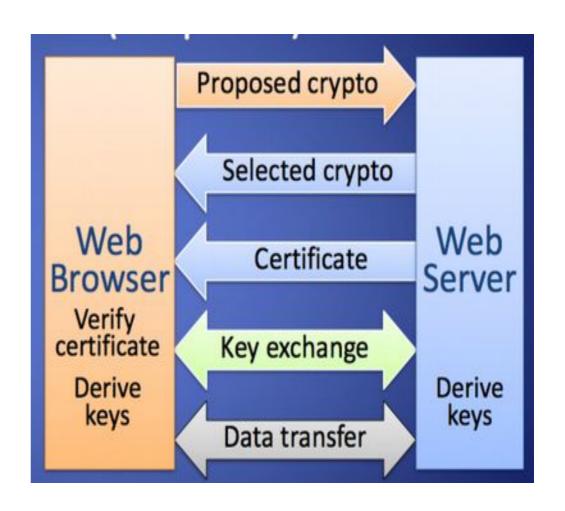


TLS: Overview

- Establish a session
 - Agree on algorithms
 - Share secrets
 - Perform authentication
- Transfer application data
 - Ensure privacy and integrity

TLS Overview

- Browser sends supported crypto algorithms
- Server picks strongest algorithms it supports
- Server sends certificate (chain)
- Client verifies certificate (chain)
- Client and server agree on secret value R by exchanging messages
- Secret value R is used to derive keys for symmetric encryption and hash-based authentication of subsequent data transfer

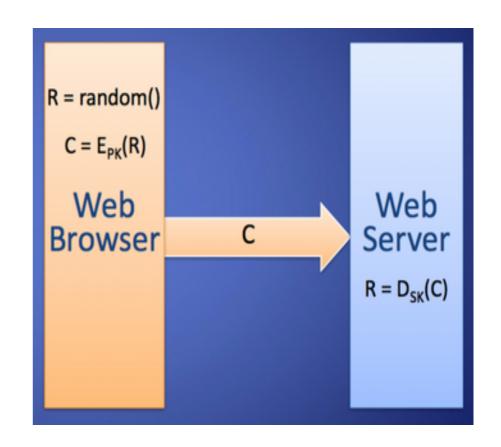


TLS:Key Exchange

- Need secure method to exchange secret key
- Use public key encryption for this
 - "key pair" is used either one can encrypt and then the other can decrypt
 - slower than conventional cryptography
 - share one key, keep the other private
- Choices are RSA or Diffie-Hellman

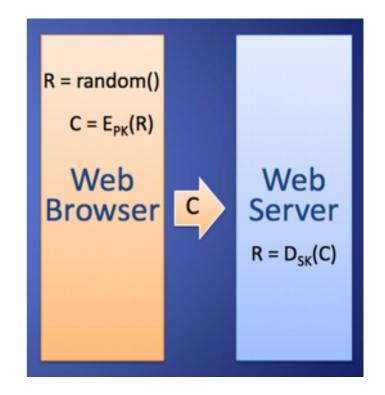
Basic Key Exchange

- Called RSA key exchange for historical reasons
- Client generates random secret value R
- Client encrypts R with public key, PK, of server C = EPK(R)
- Client sends C to server
- Server decrypts C with private key, SK, of server R = DSK(C)

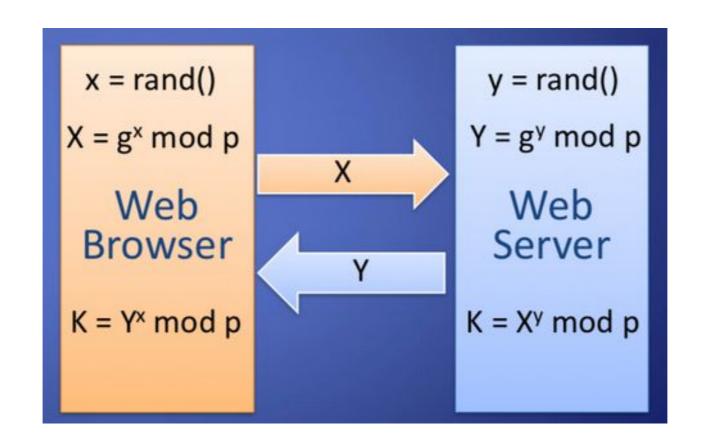


Forward Secrecy

- Compromise of public-key encryption private keys does not break confidentiality of past messages
- TLS with basic key exchange does not provide forward secrecy
- Attacker eavesdrop and stores communication
- If server's private key is compromised, attacker finds secret value R in key exchange and derives encryption keys

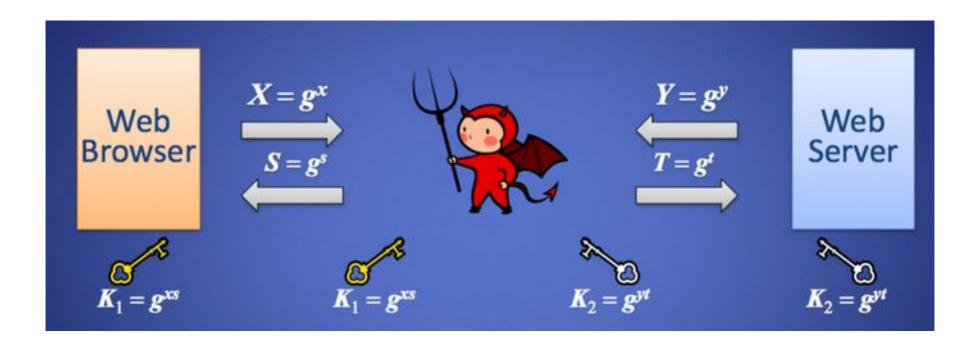


Diffie Hellman Key Exchange



Achieves forward secrecy

Attacker in the Middle



Solution:

Browser and server send signed X and Y respectively Requires each to know the public key of the other

TLS: Privacy

- Encrypt message so it cannot be read
- Use conventional cryptography with shared key
 - DES, 3DES, AES
 - RC2, RC4
 - IDEA



TLS Encrypts

- ALL Browser-Server and Server-Browser except which-browser is talking to whichserver
- URL of requested document
- Contents of requested document
- Contents of any submitted form fill-outs
- Cookies sent from browser to server
- Cookies sent from server to browser
- Contents of HTTP header
- Javascript communications
- Etc.

TLS: Integrity

- Compute fixed-length Message Authentication Code (MAC)
 - Includes hash of message
 - Includes a shared secret
 - Include sequence number
- Transmit MAC with message

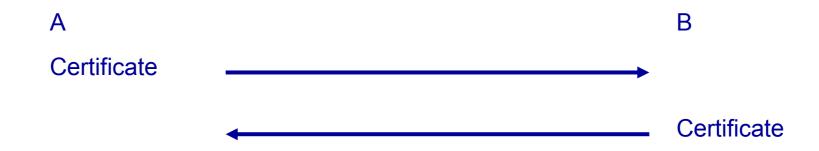
TLS: Integrity

- Receiver creates new MAC
 - should match transmitted MAC
- TLS allows MD5, SHA-1

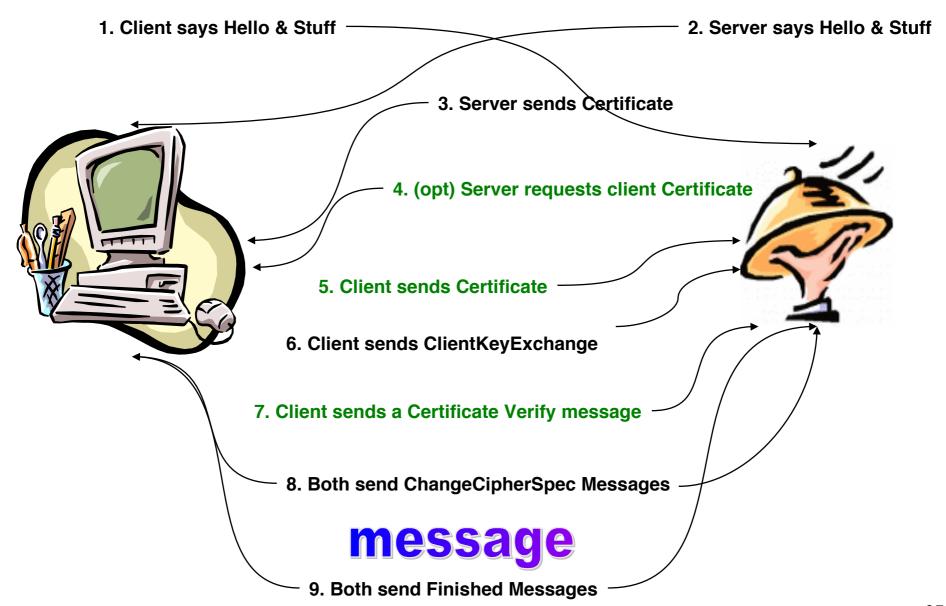


TLS: Authentication

- Verify identities of participants
- Client authentication is optional
- Certificate is used to associate identity with public key and other attributes



TLS Transaction

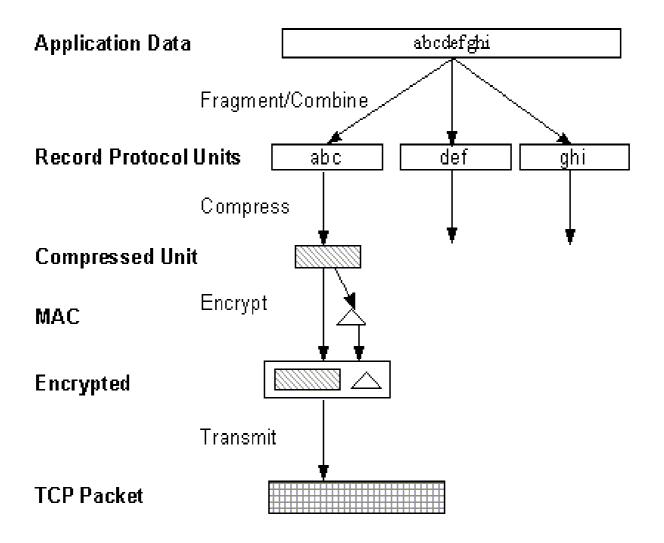


TLS: Architecture

- TLS defines Record Protocol to transfer application and TLS information
- A session is established using a Handshake Protocol

Handshake	Change	Alert		
Protocol	Cipher Spec	Protocol		
TLS Record Protocol				

TLS: Record Protocol



TLS: HTTP Application

- HTTP most common TLS application
 https://
- Requires TLS-capable web server
- Requires TLS-capable web browser

Public Key Certificates

- X.509 Certificate associates public key with identity
- Certification Authority (CA) creates certificate
 - Adheres to policies and verifies identity
 - Signs certificate
- User of Certificate must ensure it is valid

Subject Names

- X.500 Distinguished Name (DN)
- Associated with node in hierarchical directory (X.500)
- Each node has Relative Distinguished Name (RDN)
 - Path for parent node
 - Unique set of attribute/value pairs for this node

Example Subject Name

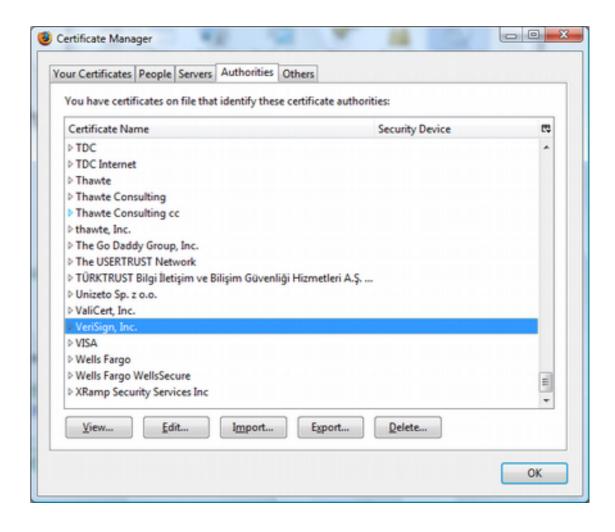
- Country at Highest Level (e.g. US)
- Organization typically at next level (e.g. CertCo)
- Individual below (e.g. Common Name "Kasun" with Id = 1)

```
DN = {
```

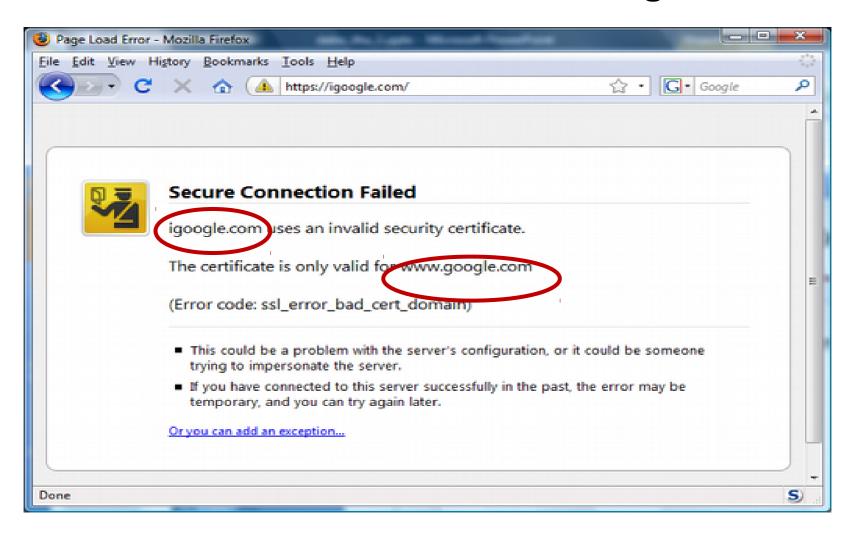
- C=LK;
- O=UCSC;
- CN=Kasun, ID=1}

Certificate Authorities

Browsers accept
certificates from a
large number of CAs



Firefox: Invalid cert dialog



Firefox 3.0: Four clicks to get firefox to accept cert

page is displayed with full HTTPS indicators

SSL Indicators

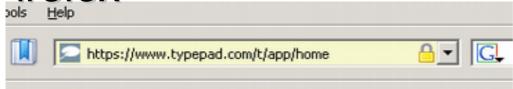
Microsoft IE



Mozilla



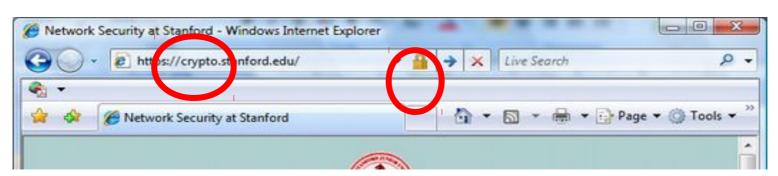
Firefox



Safari



The lock icon: SSL indicator





Intended goal:

- Provide user with identity of page origin
- Indicate to user that page contents were not viewed or modified by a network attacker

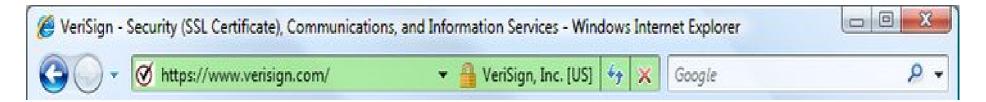
In reality:

- Origin ID is not always helpful
- Many other problems

Version 3 Certificates

- Version 3 X.509 Certificates support alternative name formats as extensions
 - X.500 names
 - Internet domain names
 - e-mail addresses
 - URLs
- Certificate may include more than one name

Extended validation (EV) certs

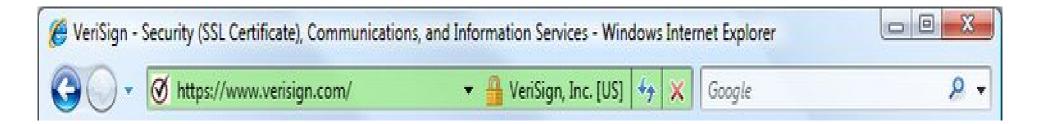


An Extended Validation Certificate (EV) is an X.509 public key certificate issued according to a specific set of identity verification criteria. These criteria require extensive verification of the requesting entity's identity by the certificate authority (CA) before a certificate is issued.

Certificates issued by a CA under the EV guidelines are not structurally different from other certificates (and hence provide no stronger cryptography than other, cheaper certificates),

Extended Validation (EV) Certs

- · Harder to obtain than regular certs
 - · requires human lawyer at CA to approve cert request
- Designed for banks and large e-commerce sites



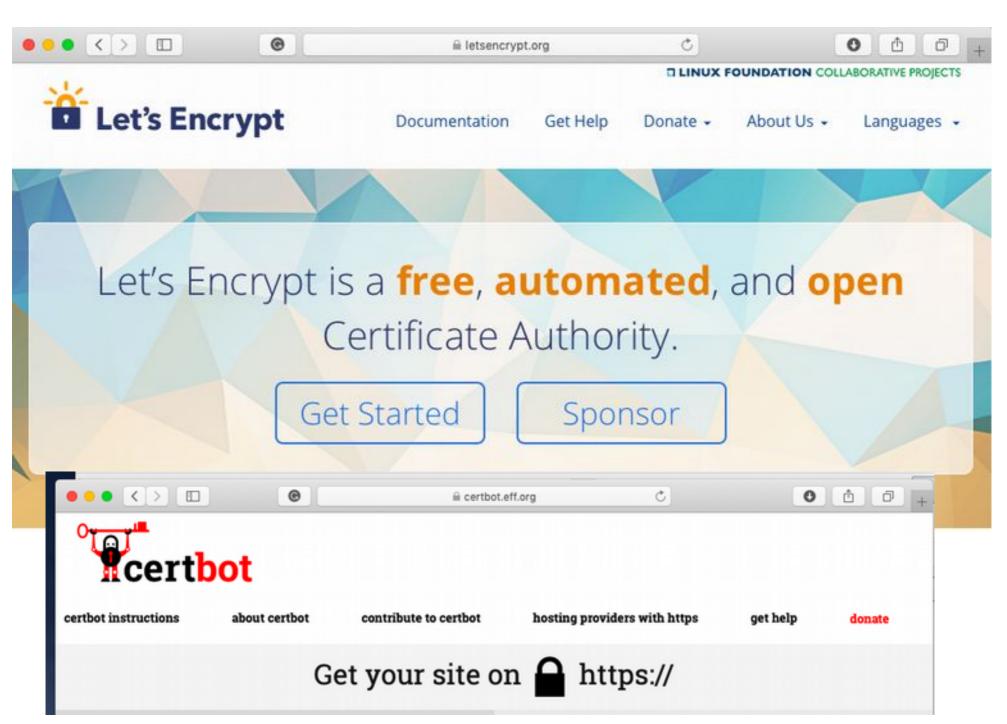
· Helps block "semantic attacks": www.bankofthevvest.com

Automatic Certificate Management Environment (ACME)

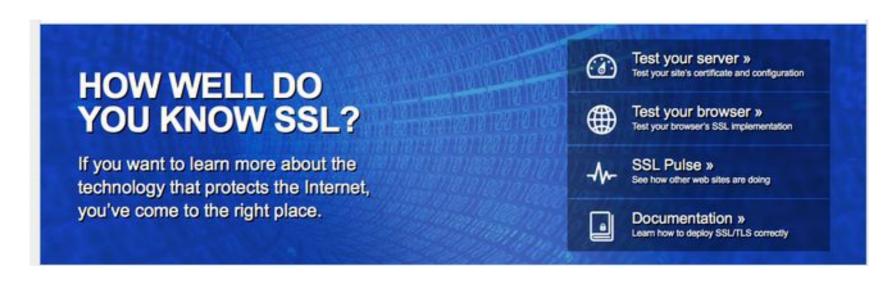
Certificates in PKI using X.509 (PKIX) are used for a number of purposes, the most significant of which is the authentication of domain names.

Thus, certificate authorities in the Web PKI are trusted to verify that an applicant for a certificate legitimately represents the domain name(s) in the certificate. Today, thi verification is done through a collection of ad hoc mechanisms.

ACME protocol automates process of verification and certificate issuance.

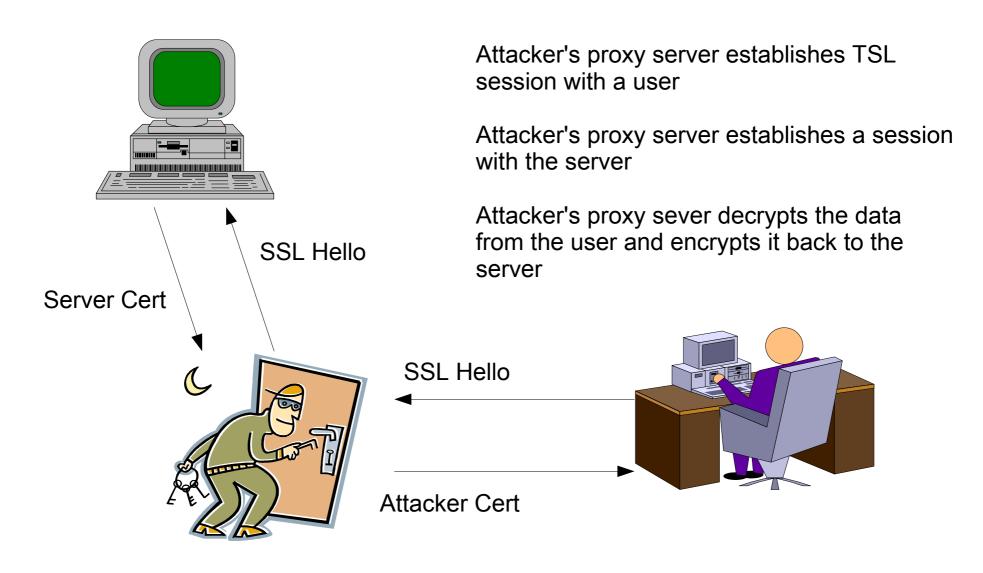


SSLABS – www.ssllabs.com

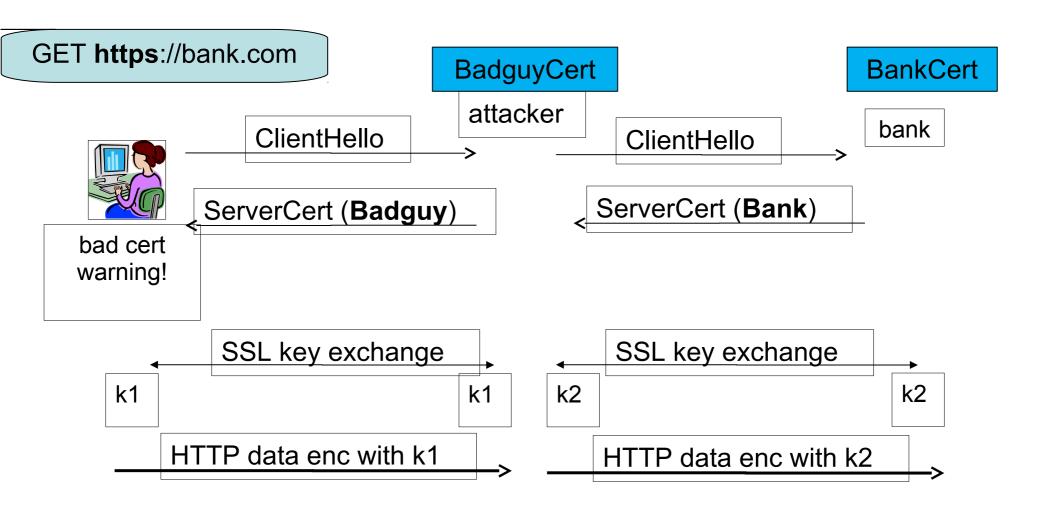




Man in the Middle



Man in the middle attack using invalid certs



Attacker proxies data between user and bank. Sees all traffic and can modify data at will.

Discussion

